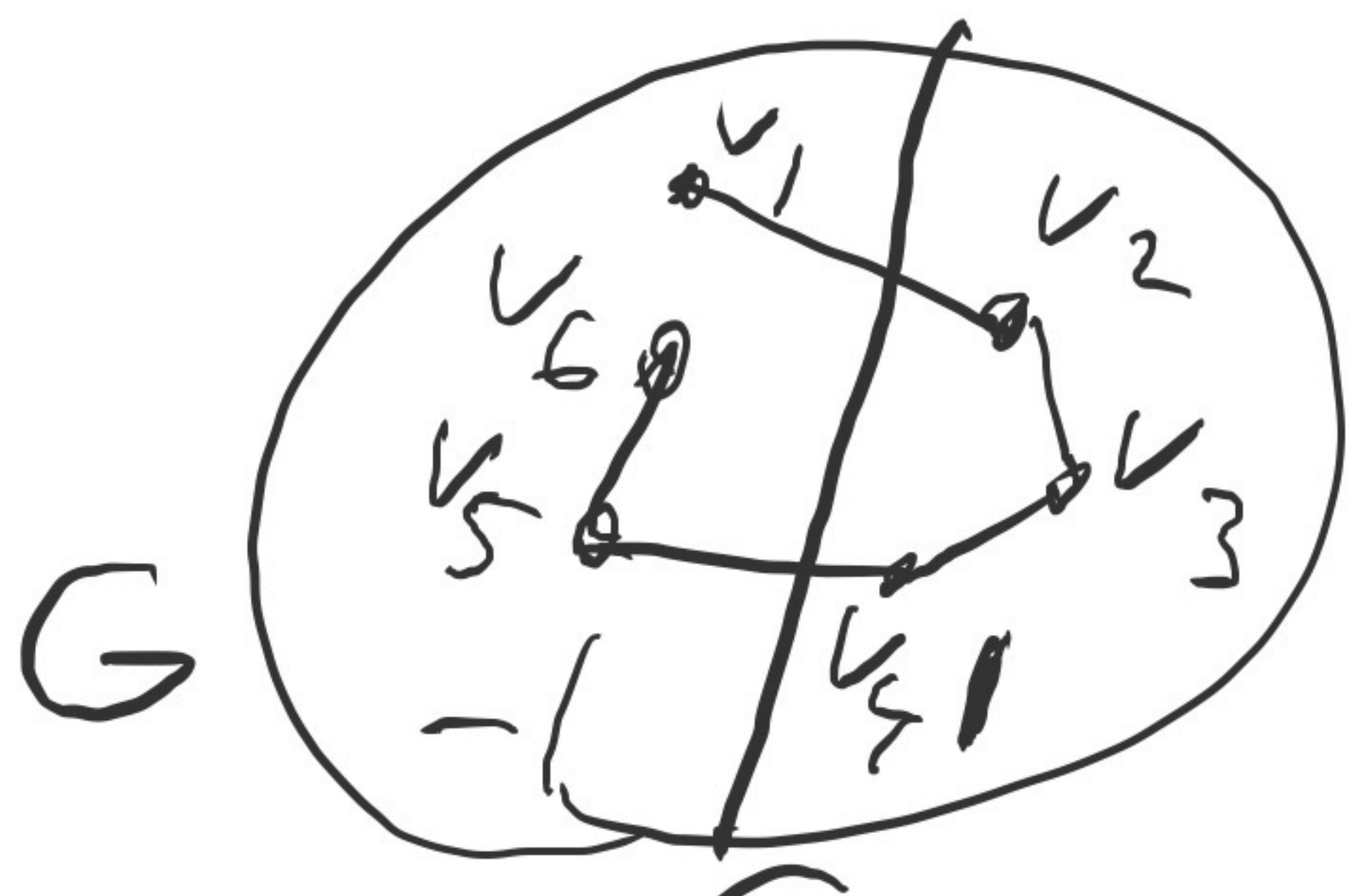


Expander Random Walks: A Fourier Analytic Approach

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Overall question: Let's say we have a d -regular expander G and we take a pseudo-random sequence as follows:

1. Label half of the vertices of G -1 and label the other half 1
2. Take a $(t-1)$ step random walk v_1, v_2, \dots, v_t and output the labels of the vertices



$-1, 1, 1, 1, -1, -1$

Q: For which functions does this behave like a random $x \in \{-1, 1\}^t$?

This paper uses Fourier analysis to show that

Symmetric functions, functions in AC^0 , and functions computable by a low width read once branching program are fooled.

Preliminaries: λ -spectral expanders

For a d -regular graph G , we'll also use G to denote the normalized adjacency matrix.

$$G_{ij} = \begin{cases} \frac{1}{d} & \text{if } (i,j) \in E(G) \\ 0 & \text{otherwise} \end{cases}$$

Definition: We say that G is a λ -spectral expander if for all eigenvectors except for $\vec{1}$, the magnitude of its eigenvalue is at most λ .

Definition: Define J to be $\frac{1}{n}$ (all ones matrix) $\forall i,j \ J_{ij} = \frac{1}{n}$

Key property: $G = J + E$ where $\|E\| \leq \lambda$ and $EJ = JE = 0$.

Preliminaries: Fourier Analysis Over $\{-1, 1\}^+$

Fourier characters: $X_S = \prod_{j \in S} x_j$

Fourier coefficients: $\hat{f}_S = E_{x \in \{-1, 1\}^+} [f(x) X_S(x)]$

Fourier decomposition: $f = \sum_{S \subseteq [t]} \hat{f}_S X_S$

Parseval's Theorem: $\sum_{S \subseteq [t]} \hat{f}_S^2 = E_{x \in \{-1, 1\}^+} [f^2]$

We'll be analyzing ± 1 -valued functions
so this will be \mathbb{I}_0 .

Preliminaries: Difference From Random

Def: Given a graph G and a labeling $\text{val}: V(G) \rightarrow \{-1, 1\}$
(where half of the vertices are labeled 1)

define $RW_{G, \text{val}, t}$ to be the distribution obtained by taking the labels of a $(t-1)$ step random walk

Definition: Given a function $f: \{-1, 1\}^t \rightarrow \{-1, 1\}$,

define $\epsilon_{G, \text{val}}(f) = \mathbb{E}_{x \sim RW_{G, \text{val}, t}} [f(x)] - \mathbb{E}_{x \sim \{-1, 1\}^t} [f(x)]$

Define $\epsilon_\lambda(f) = \max_{\text{val } \lambda\text{-spectral expanders } G} |\epsilon_{G, \text{val}}(f)|$

Idea: Bound $\epsilon_\lambda(f)$ by analyzing $\epsilon_{G, \text{val}}(X_5)$

Results on Symmetric Functions

Definition: We say that a function

$f: \{-1, 1\}^t \rightarrow \{-1, 1\}$ is symmetric if

$f(x_1, x_2, \dots, x_t)$ only depends on the number of -1 s and 1 s in x_1, \dots, x_t .

Theorem 1.1': For all symmetric functions f ,

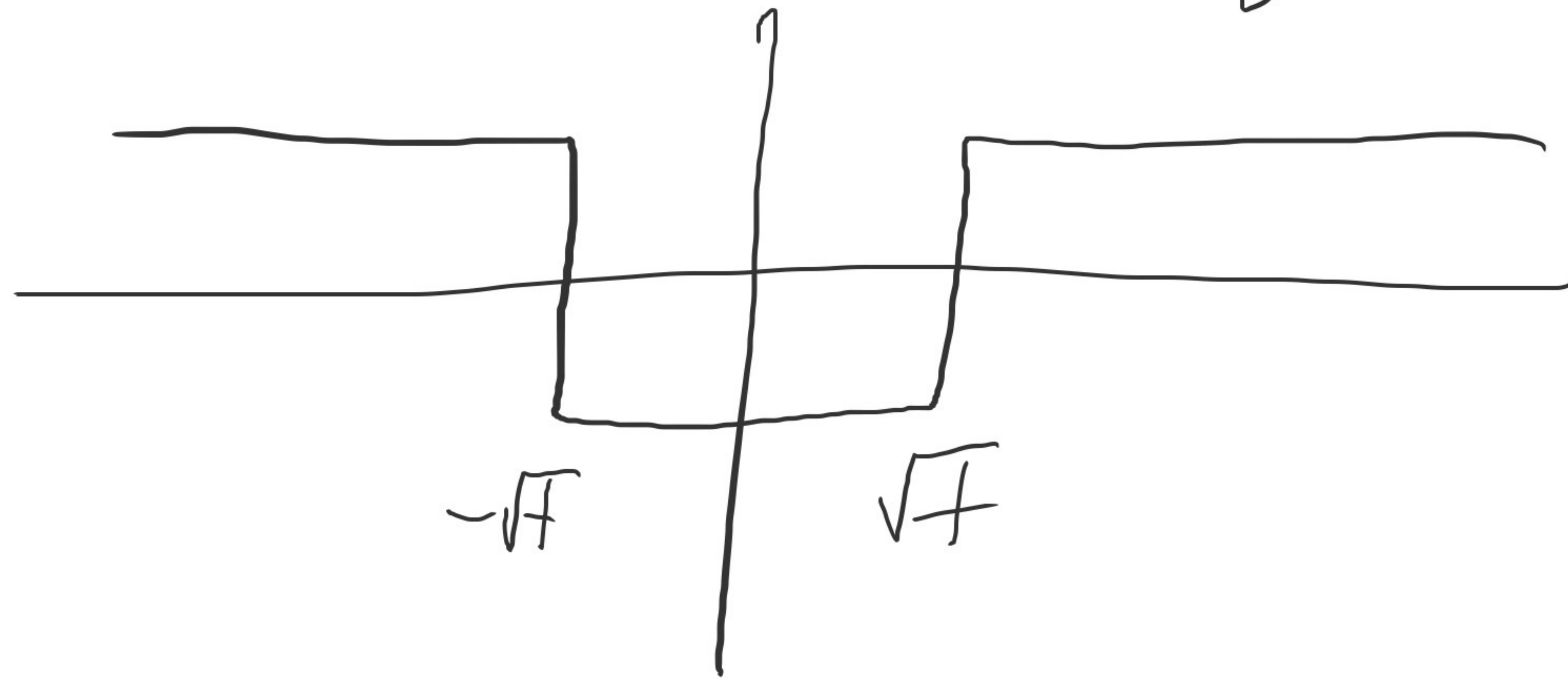
$$\Sigma_\lambda(f) \text{ is } O(\lambda).$$

Restatement: If we only consider the number of 1 s and -1 s in x , the total variation distance between $RW_{G, \text{val}, t}$ and $U_t = \text{uniform dist on } \{-1, 1\}^t$ is $O(\lambda)$

Theorem 1.4:

$$\Sigma_\lambda(\text{MAJ}_t) \text{ is } O\left(\frac{\lambda^2}{\sqrt{t}}\right)$$

I think $O(\lambda)$ is tight. I think



is an example.

Note: $RW_{G, val, t}$ does not necessarily fool all functions



If $\Pr(\text{crossing}) = \frac{1}{2} + \Omega(\lambda)$ then

for $RW_{G, val, t}$, $E[\# \text{ of crossings}]$

is $\frac{t}{2} + \Omega(\lambda t)$
 For U_t , $\# \text{ of crossings}$ is $\frac{t}{2} \pm o(\sqrt{t})$.

Results on General Functions

Definition: Given a function $f: \{-1, 1\}^t \rightarrow \{-1, 1\}$,
define $L_{1,k}(f) = \sum_{\substack{S \subseteq [t] \\ |S|=k}} |\hat{f}_S|$

Definition: Define $L_1^+(b)$ to be the family
of functions $f: \{-1, 1\}^t \rightarrow \{-1, 1\}$
such that

$$\forall k \quad L_{1,k}(f) \leq b^k$$

Claim 5.2': For any function $f \in L_1^+(b)$,
 $\epsilon_\lambda(f)$ is $O(\lambda b^2)$

Avishay Tal

Tight Bounds on the Fourier spectrum of AC^0

Theorem: Any function computable by a size s depth d AC^0 circuit is in $L_1^+(b)$ for $b = O((\log s)^{d-1})$

$$\sum_{\lambda} \hat{f}(\lambda) \text{ is } O(\lambda (\log s)^{2(d-1)})$$

Eshan Chattopadhyay, Pooya Hatami, Omer Reingold, and Avishay Tal

Improved Pseudorandomness for Unordered Branching Programs Through Local Monotonicity

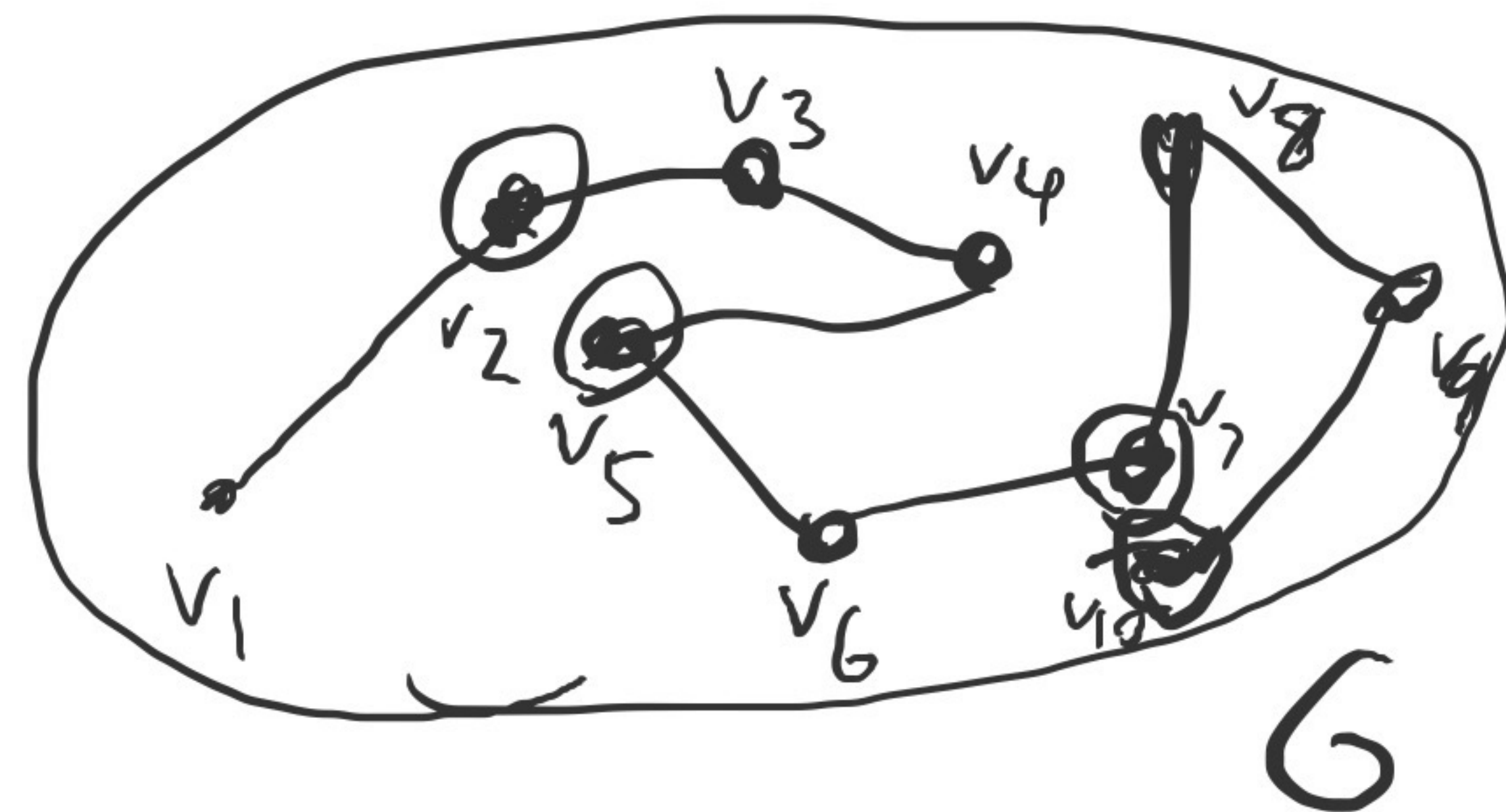
Theorem: If f can be computed by a width w read once branching program, $b = O((\log t)^w)$

$$\sum_{\lambda} \hat{f}(\lambda) \text{ is } O(\lambda (\log t)^{2w})$$

$$b \leq DT(f) \quad \Sigma_\lambda(f) \text{ is } O(\lambda DT(f)^2).$$

Analysis:

$$\begin{aligned} \text{Consider } \Sigma_{G, \text{val}}(X_S) \\ = \underset{RW_{G, \text{val}}}{E} [X_S(x)] \end{aligned}$$



Define P to be the diagonal matrix where $P_{ii} = \text{val}(i)$.

$$S = \{2, 5, 7, 10\}$$

$$\Delta_0 = 1 \quad \Delta_1 = 3 \quad \Delta_2 = 2 \quad \Delta_3 = 3$$

$$\Sigma_{G, \text{val}}(X_S) = \underset{S}{\mathbb{1}}^T (PG^{\Delta_1}) \dots (PG^{\Delta_1})(PG^{\Delta_0}) \underset{S}{\mathbb{1}}$$

Idea: We'll analyze this by breaking G^{Δ_j} up into $G^{\Delta_j} = J + E^{\Delta_j}$ (recall $G = J + E$ where $\|E\| \leq \lambda$)

$$\Sigma_{\text{Gral}}(X_S) = \vec{1}^T (PG^{\Delta_{k-1}}) \dots (PG^{\Delta_2})(PG^{\Delta_1}) P \frac{\vec{1}}{n}$$

Idea: Consider which E^{Δ_j} terms we have.

Proposition: $\vec{1}^T P \vec{1} = 0$. Since $J = \frac{1}{n} \vec{1} \cdot \vec{1}^T$,

$$JP\vec{1} = 0 \quad \vec{1}^T PJ = 0 \quad JPJ = 0.$$

Claim: If we don't have E^{Δ_1} or $E^{\Delta_{k-1}}$, this part is 0.

If we don't have E^{Δ_j} and $E^{\Delta_{j+1}}$ for some j , this part is 0.

Definition: Define \mathcal{F}_k to be the set of $I \subseteq [k-1]$ such that

1. $1 \in I, k-1 \in I$,
2. $\forall j \in [k-2] \quad j \in I \text{ or } j+1 \in I$.

$$\text{Proposition 4.2: } \left| \Sigma_{\text{Gral}}(X_S) \right| \leq \sum_{I \in \mathcal{F}_k} \lambda^{\sum_{j \in I} \Delta_j(S)}$$

Proposition 4.2: $|\Sigma_{G, \text{val}}(X_S)| \leq \sum_{I \in \mathcal{F}_k} \lambda^{\sum_{i \in I} \Delta_i(S)}$

Examples: $|S|=2 \leq \lambda^{\Delta_1} \quad \mathcal{F}_2 = \{\{1,3\}\}$

$|S|=3 \leq \lambda^{\Delta_1 + \Delta_2} \quad \mathcal{F}_3 = \{\{1,2,3\}\}$

$|S|=4 \leq \lambda^{\Delta_1 + \Delta_3} + \lambda^{\Delta_1 + \Delta_2 + \Delta_3} \quad \mathcal{F}_4 = \{\{1,3\}, \{1,2,3\}\}$

$|S|=6 \quad \mathcal{F}_6 = \{\{1,3,5\}, \{1,2,4,5\}, \{1,2,3,5\}, \{1,3,4,5\}, \{1,2,3,4,5\}\}$

Proposition: $\forall I \in \mathcal{F}_k, |I| \geq \lceil \frac{k}{2} \rceil$

Proposition: If $|S|=0$ or $1 \quad \Sigma_{G, \text{val}}(X_S) = 0$

Corollary: $\forall S: |S| \geq 2, |\Sigma_{G, \text{val}}(X_S)| \leq 2^k \cdot \lambda^{\lceil \frac{k}{2} \rceil} \leq (4\lambda)^{\frac{k}{2}}$

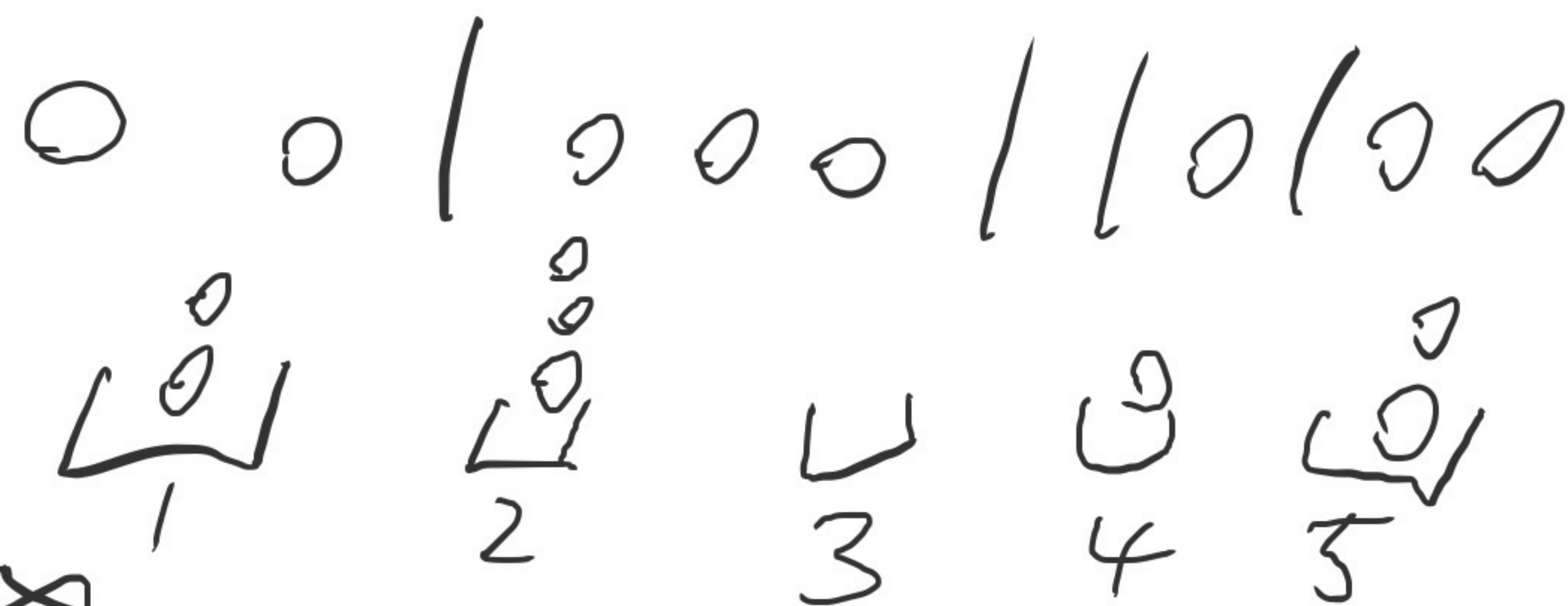
Corollary: If $f \in L_1^+(b)$ then

$\Sigma_\lambda(f) \leq \sum_{k=2}^{\infty} (4\lambda)^{\frac{k}{2}} \cdot b^k$ which is $O(\lambda b^2)$.

Review: Unlabeled Balls in Labeled Bins

Proposition: There are $\binom{n+k-1}{k-1}$ ways to put n unlabeled balls in k labeled bins.

Trick: Imagine putting $k-1$ dividing lines among the n balls



Claim 3.1: $\sum_{j=r}^{\infty} \binom{j}{r} \lambda^j = \frac{\lambda^r}{(1-\lambda)^{r+1}}$ if $|\lambda| < 1$

Proof: $\frac{1}{1-\lambda} = 1 + \lambda + \lambda^2 + \lambda^3 + \dots$

What is the coefficient of λ^j in $\frac{\lambda^r}{(1-\lambda)^{r+1}}$? Claim: This is like putting $j-r$ balls in $r+1$ labeled bins. There are $\binom{j-r+r+1-1}{r+1-1} = \binom{j}{r}$ ways to do this.

Definition: Define $\beta_k = \sum_{\substack{S \subseteq [t]: \\ |S|=k}} \sum_{\text{bval}} (x_S)$

Lemma 4.4: $|\beta_k| \leq 2^k \binom{t-1}{\lfloor \frac{k}{2} \rfloor} \left(\frac{\lambda}{1-\lambda}\right)^{\lfloor \frac{k}{2} \rfloor}$

Claim 4.5: $|\beta_k| \leq 2^k \sum_{m=\lfloor \frac{k}{2} \rfloor}^{t-1} \binom{m-1}{\lfloor \frac{k}{2} \rfloor - 1} \binom{t-m}{k - \lfloor \frac{k}{2} \rfloor} \lambda^m$

By symmetry, this only depends on $|I|$ and is decreasing in I .

Proof: $|\beta_k| \leq \sum_{\substack{S \subseteq [t] \\ |S|=k}} \sum_{I \in \mathcal{F}_k} \lambda^{\sum_{j \in I} \Delta_j(S)} = \sum_{I \in \mathcal{F}_k} \sum_{\substack{S \subseteq [t] \\ |S|=k}} \lambda^{\sum_{j \in I} \Delta_j(S)}$

Choosing a particular $I^* \in \mathcal{F}_k$ where $|I^*| = \lfloor \frac{k}{2} \rfloor$, $|\beta_k| \leq 2^k \sum_{\substack{S \subseteq [t] \\ |S|=k}} \lambda^{\sum_{j \in I^*} \Delta_j(S)}$

Q: What is the coefficient of λ^m in this expression?

We need to split m factors of λ among $|I^*|$ Δ_j where each Δ_j is at least 1. This is = putting $m - |I^*|$ balls in $|I^*|$ labeled bins.

We also need to divide the remaining $t-1-m$ steps of the path among $k+1 - |I^*|$ other Δ . The middle Δ must be at least 1.

 This is = putting $t-1-m - (k+1 - |I^*| - 2)$ unlabeled balls in $k+1 - |I^*|$ labeled bins.

This gives a coefficient of $\binom{m-1}{|I^*|-1} \binom{t-m}{k-|I^*|}$ for λ^m

Claim 4.5: $|\beta_k| \leq 2^k \sum_{m=\lceil \frac{k}{2} \rceil}^{+\infty} \binom{m-1}{\lceil \frac{k}{2} \rceil - 1} \binom{t-m}{k - \lceil \frac{k}{2} \rceil} \lambda^m$

$$\leq 2^k \binom{t-1}{k - \lceil \frac{k}{2} \rceil} \lambda \sum_{m=\lceil \frac{k}{2} \rceil}^{\infty} \binom{m-1}{\lceil \frac{k}{2} \rceil - 1} \lambda^{m-1}$$

Claim 3.1: $\sum_{j=r}^{\infty} \binom{j}{r} \lambda^j = \frac{\lambda^r}{(1-\lambda)^{r+1}} \quad \text{or} \quad \frac{\lambda^{\lceil \frac{k}{2} \rceil - 1}}{(1-\lambda)^{\lceil \frac{k}{2} \rceil}}$

$$r = \lceil \frac{k}{2} \rceil - 1$$

Lemma 4.4

$$|\beta_k| \leq 2^k \binom{t-1}{k - \lceil \frac{k}{2} \rceil} \cdot \frac{\lambda^{\lceil \frac{k}{2} \rceil}}{(1-\lambda)^{\lceil \frac{k}{2} \rceil}}$$

Lemma 4.4: $|\beta_k| \leq 2^k \binom{+1}{\lfloor \frac{k}{2} \rfloor} \left(\frac{\lambda}{1-\lambda}\right)^{\lceil \frac{k}{2} \rceil}$

For a symmetric function f , defining \hat{f}_k to be \hat{f}_S for some S of size k

$$\Sigma_\lambda(f) \leq \sum_{k=2}^+ |\beta_k| \cdot \hat{f}_k$$

How large can \hat{f}_k be?

There are $\binom{+}{k}$ subsets S of size k .

Since $\sum_{S: |S|=k} \hat{f}_S^2 = \binom{+}{k} \hat{f}_k^2 \leq 1$ $\hat{f}_k \leq \frac{1}{\sqrt{\binom{+}{k}}}$

$$\Sigma_\lambda(f) \leq \sum_{k=2}^+ 2^k \frac{\binom{+1}{\lfloor \frac{k}{2} \rfloor}}{\sqrt{\binom{+}{k}}} \left(\frac{\lambda}{1-\lambda}\right)^{\lceil \frac{k}{2} \rceil}$$

$O\left(\frac{\lambda^{\lceil \frac{k}{2} \rceil}}{\sqrt{k}}\right)$ k is odd

$O\left(\lambda^{\frac{k}{2}}\right)$ k is even