

# Local Analysis of Higher-Order Random Walks and Applications

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## 1 Tighter Spectral Analysis

These notes are based on Vedat Levi Alev exposition in the HDX-Codes cluster.

Define the operators  $P_k^\nabla := U_{k-1}D_k$  and  $P_k^\Delta := D_{k+1}U_k$  where  $U_k: \mathbb{R}^{X(k)} \rightarrow \mathbb{R}^{X(k+1)}$  and  $D_k: \mathbb{R}^{X(k)} \rightarrow \mathbb{R}^{X(k-1)}$  are the usual up and down operators of a weighted simplicial complex  $X(\leq d)$ . For convenience, we will be working with size rather than dimension.

**Theorem 1.1** (Kaufman-Oppenheim [3]).

$$\lambda_2(P_k^\nabla) \leq 1 - \frac{1}{k+1} + \frac{k}{2} \cdot \gamma,$$

where  $\gamma = \max_{S \in X(\leq k-2)} \lambda_2(A_S)$ .

**Remark 1.2.** *The above result requires the links to be  $\gamma = O(1/k^2)$  expanding to give a non-trivial bound on  $\lambda_2(P_k^\nabla)$ .*

Alev et al. show that the  $P_k^\nabla$  walk satisfy  $\lambda_2(P_k^\nabla) < 1$  provided  $\gamma < 1$ . More precisely, they showed the following.

**Theorem 1.3** (Alev et al. [1]). *For  $k \in [d]$ , we have*

$$\lambda_2(P_{k-1}^\Delta) = \lambda_2(P_k^\nabla) \leq 1 - \frac{1}{k} \prod_{i=0}^{k-2} (1 - \gamma_i),$$

where  $\gamma_i = \max_{S \in X(i)} \lambda_2(A_S)$  and  $A_S$  is the normalized adjacency operator of the link  $X_S$ .

Define and  $P_k^\wedge := (k+1)/k \cdot (D_k U_{k+1} - I/(k+1))$ , i.e.,  $P_k^\wedge$  is the non-lazy version of  $P_k^\Delta$ . The key lemma is the following.

**Lemma 1.4** (Key Lemma). *For  $k \in [d-1]$ , we have*

$$P_k^\wedge - P_k^\nabla \preceq \gamma_{k-1}(I - P_k^\nabla).$$

A version of Lemma 1.4 was known with RHS  $\gamma_{k-1}I$ . The change to  $\gamma_{k-1}(I - P_k^\nabla)$  was crucial to the improved bound on  $\lambda_2(P_k^\nabla)$ . We omit the proof of Lemma 1.4, but observe that it follows from a finer account of the mass of “parallel vectors on the links” similar in spirit to the proof of Oppenheim Trickleing-down Theorem.

*Proof of Theorem 1.3.* The equality  $\lambda_2(P_{k-1}^\Delta) = \lambda_2(P_k^\nabla)$  readily follows from  $P_{k-1}^\Delta := D_k U_{k-1}$  and  $P_k^\nabla := U_{k-1} D_k$  having the same multiset of non-zero eigenvalues.

Now, the proof proceeds by induction. Note that  $P_1^\nabla = U_0 D_1$  is rank one and thus  $\lambda_2(P_1^\nabla) = 0$  proving the base case.

From Lemma 1.4, we have  $P_k^\wedge \preceq \gamma_{k-1} I + (1 - \gamma_{k-1}) P_k^\nabla$  implying

$$\begin{aligned} \lambda_2(P_k^\wedge) &\leq \gamma_{k-1} + (1 - \gamma_{k-1}) \lambda(P_k^\nabla) \\ &\leq \gamma_{k-1} + (1 - \gamma_{k-1}) \left(1 - \frac{1}{k} \prod_{i=0}^{k-2} (1 - \gamma_i)\right) \\ &= 1 - \frac{1}{k} \prod_{i=0}^{k-1} (1 - \gamma_i). \end{aligned}$$

Since  $P_k^\wedge := (k+1)/k \cdot (P_k^\Delta - I/(k+1))$ , we concluded that

$$\lambda_2(P_k^\Delta) \leq 1 - \frac{1}{k+1} \prod_{i=0}^{k-1} (1 - \gamma_i),$$

also giving a bound for  $\lambda_2(P_{k+1}^\nabla) = \lambda_2(P_k^\Delta)$ . □

## 2 Applications

For instance, fast mixing of the high-dimensional walks has found applications in counting/sampling combinatorial objects [2]. Therefore, it is highly desirable to reduce the expansion requirement of the links as much as possible.

## References

- [1] Vedat Levi Alev, Nima Anari, Lap Chi Lau, Kuikui Liu, and Shayan Oveis Gharan. A note on the second eigenvalue of higher order random walks. 2019. <https://cs.uwaterloo.ca/~vlalev/notes/higher-order.pdf>.
- [2] Nima Anari, Kuikui Liu, Shayan Oveis Gharan, and Cynthia Vinzant. Log-concave polynomials II: high-dimensional walks and an FPRAS for counting bases of a matroid. *CoRR*, abs/1811.01816, 2018.
- [3] Tali Kaufman and Izhar Oppenheim. High order random walks: Beyond spectral gap. In *Approximation, Randomization, and Combinatorial Optimization. Algorithms and Techniques, APPROX/RANDOM 2018, August 20-22, 2018 - Princeton, NJ, USA*, pages 47:1–47:17, 2018.